

Chapter 2: Application layer

- ❑ 2.1 Principles of network applications
 - ❖ app architectures
 - ❖ app requirements
- ❑ 2.2 Web and HTTP
- ❑ 2.4 Electronic Mail
 - ❖ SMTP, POP3, IMAP
- ❑ 2.5 DNS
- ❑ 2.6 P2P applications
- ❑ 2.7 Socket programming with TCP
- ❑ 2.8 Socket programming with UDP

So is it Peer to Peer or P2P???

Peer to Peer

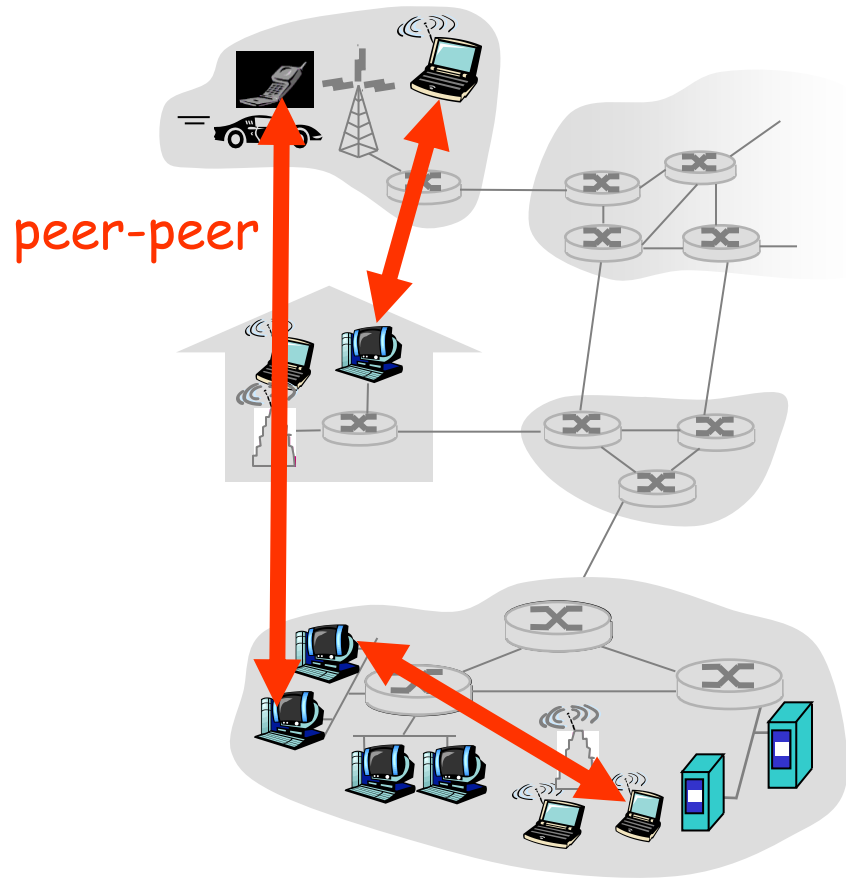
- ❑ Done typically on a LAN
- ❑ Uses software components built into the OS
- ❑ Protocol = TCP/IP
- ❑ Legally share devices and files

P2P

- ❑ Done typically on the internet
- ❑ Must install special 3rd party software (ares, trustyfiles, bearshare)
- ❑ Additional protocols (BitTorrent, Ants P2P, eDonkey)
- ❑ Pirated software, music and movies

Pure P2P architecture

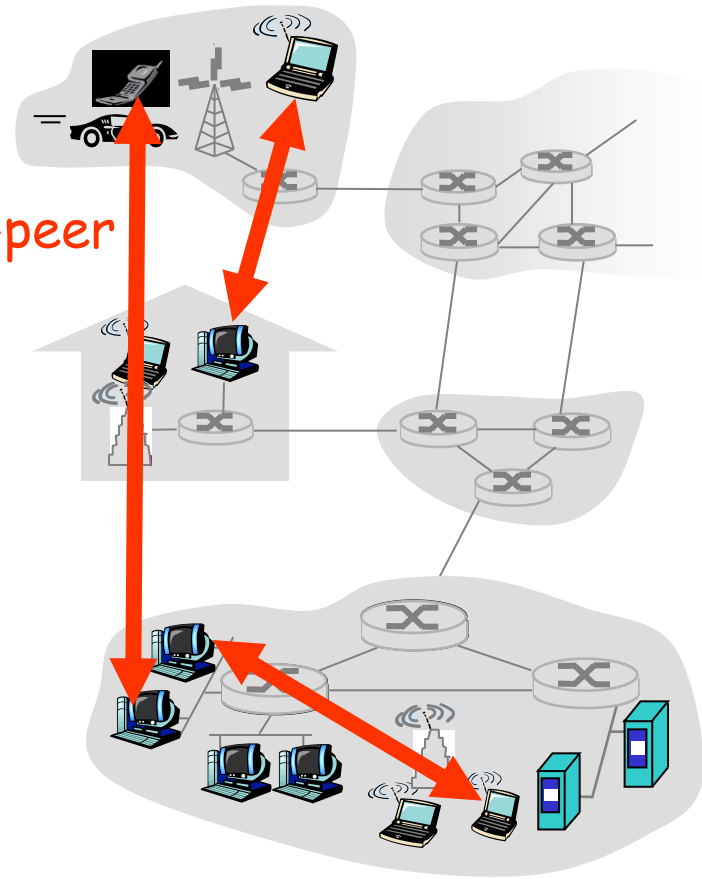
- ❑ LAN model
- ❑ OS: Mac, Vista, XP, Linux
- ❑ Wireless or Wire
- ❑ Environments: Home, Office
- ❑ GUID: Unique Identification



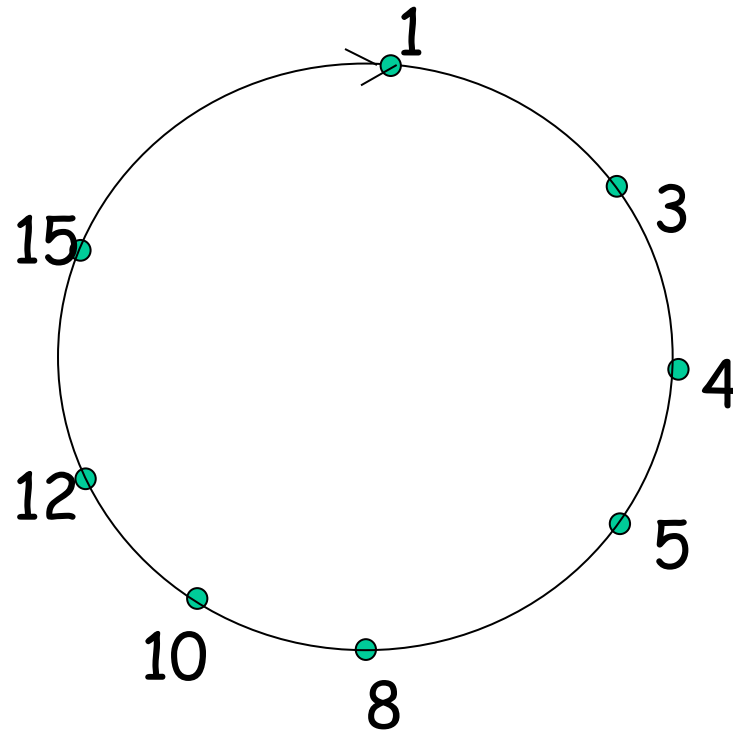
Pure P2P architecture

- ❑ *no* always-on server
- ❑ arbitrary end systems directly communicate
- ❑ peers are intermittently connected and change IP addresses

peer-peer



Peer Arrangement



- ❑ Only maintain reference to a subset of the peers (exponentially further and further away from them)

Advantages of Peer-to-Peer

Simple

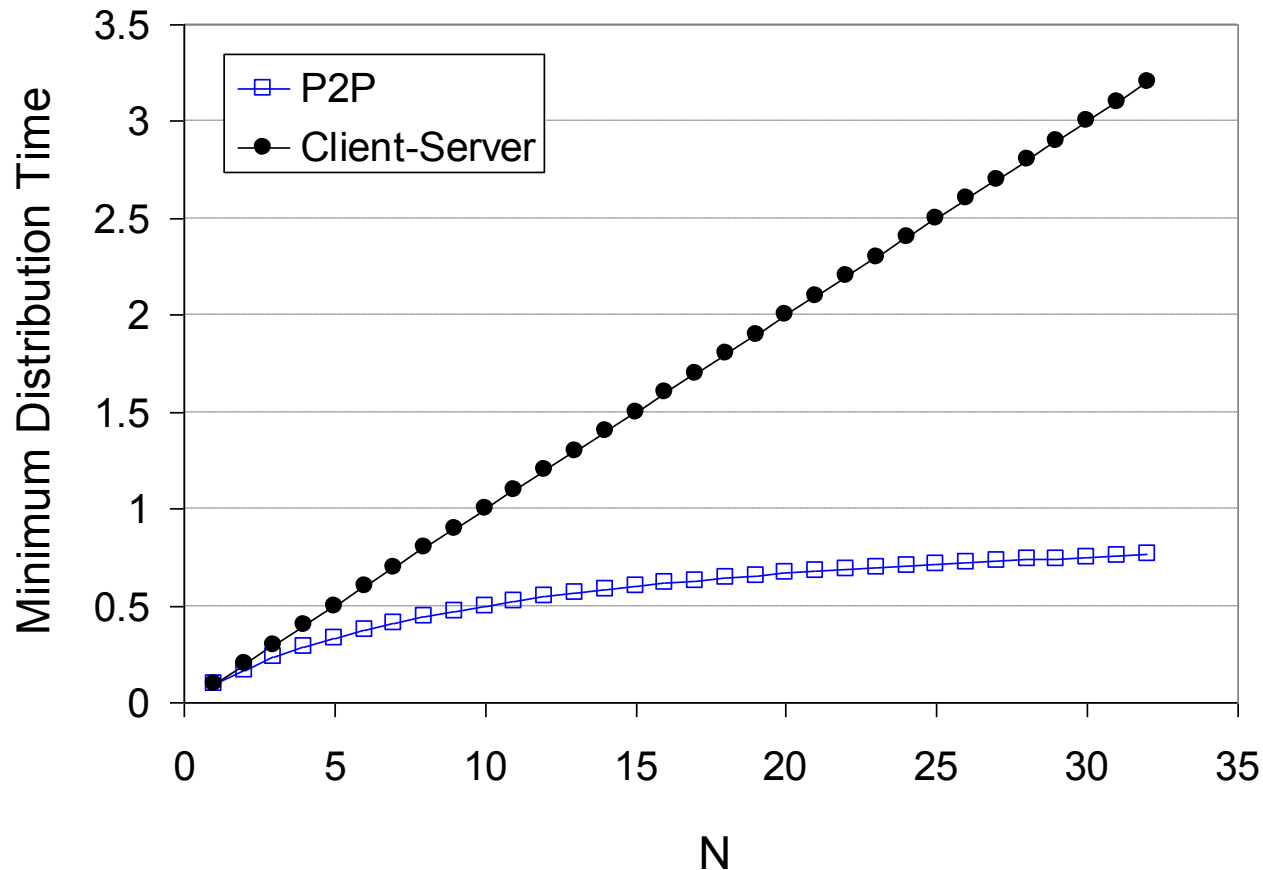
- ❑ For users to understand once setup
- ❑ Small companies do not have to hire IT staff
- ❑ Practical for small business and home offices (SOHOs)

Low Cost

- ❑ Media needed: hub or switch, NIC, cables
- ❑ Wireless technology
- ❑ Networking software comes with OS

Server-client vs. P2P: example

Client upload rate = u , $F/u = 1$ hour, $u_s = 10u$, $d_{\min} \geq u_s$



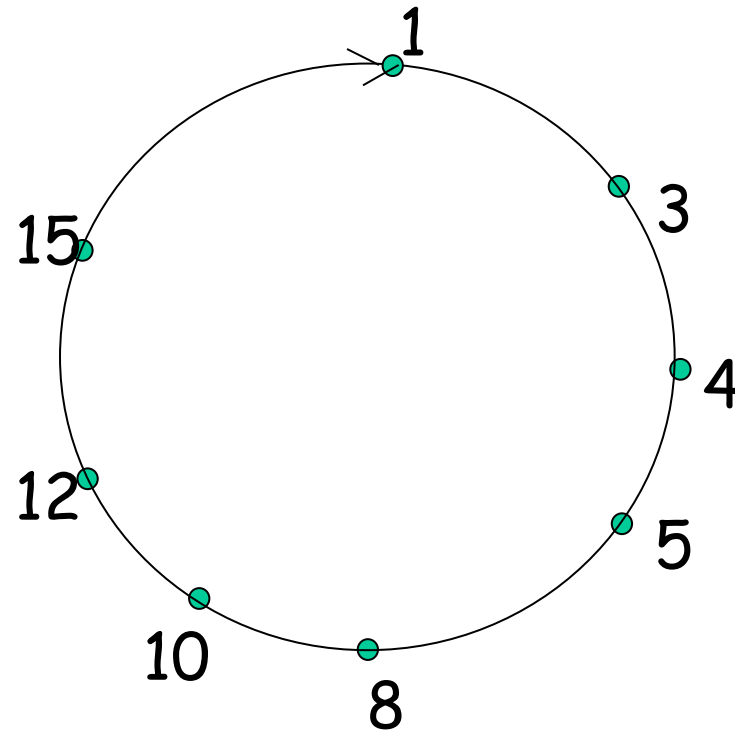
Client/ Server Based Network?

- ❑ Networking using special computers known as file servers, to process data for and facilitate communication between other computers on the network
 - ❖ File Server: manages shared resources, uses special system software designed to manage data and other resources
 - ❖ Client/Workstation: requests another computer on the network

Advantages of Client/Server

- ❑ Security
- ❑ Centralized user accounts
- ❑ Access to multiple shared resources can be centrally granted
- ❑ Optimized to handle heavy processing loads and dedicated to handling requests from clients
- ❑ Can scale to thousands of computers
- ❑ Higher license fees
- ❑ Back ups data only required only at the server

Peer to Peer Algorithms



Components for Peer-to-Peer

□ Hardware

- ❖ NIC, modem, USB port
- ❖ Cables or phone line, electrical outlet, wireless transmitter and receiver, USB cables and hub
- ❖ Wireless home routers or switch (if just two computers you need only a crossover cable)

□ Software

- ❖ Windows XP/Vista
- ❖ Linux/Mac OS with SMB services

Software Pieces to Vista/XP

- ❑ **Client:** Software installed allowing your workstation to view across the network into a device/resource/object on the network, access data on other systems e.g. Client for Microsoft networks (PCs only), client NFS (interoperability)
- ❑ **Protocol:** the language of network communication, TCP/IP
- ❑ **Services:** additional network software to provide network monitoring, QoS, remote backup, server services, virtualization etc (File and Print Service sharing)
- ❑ **NIC Driver:** Software to interact with NIC

Services

Advanced functions for a complex environment

- ❑ **Server services (file & print services):** found in all operating systems allows the sharing of files, folders, drivers, applications and printers on network
- ❑ **Backup agents:** allows a server based backup system to remotely backup the computer (*third party*)
- ❑ **Network Access Protection agent:** force clients to meet corporate policies related to required updates-patches, current anti-virus software (*third party*)
- ❑ **QoS:** Gives priority to TCP/IP packets determined by the network administrator, VoIP

What can we share?

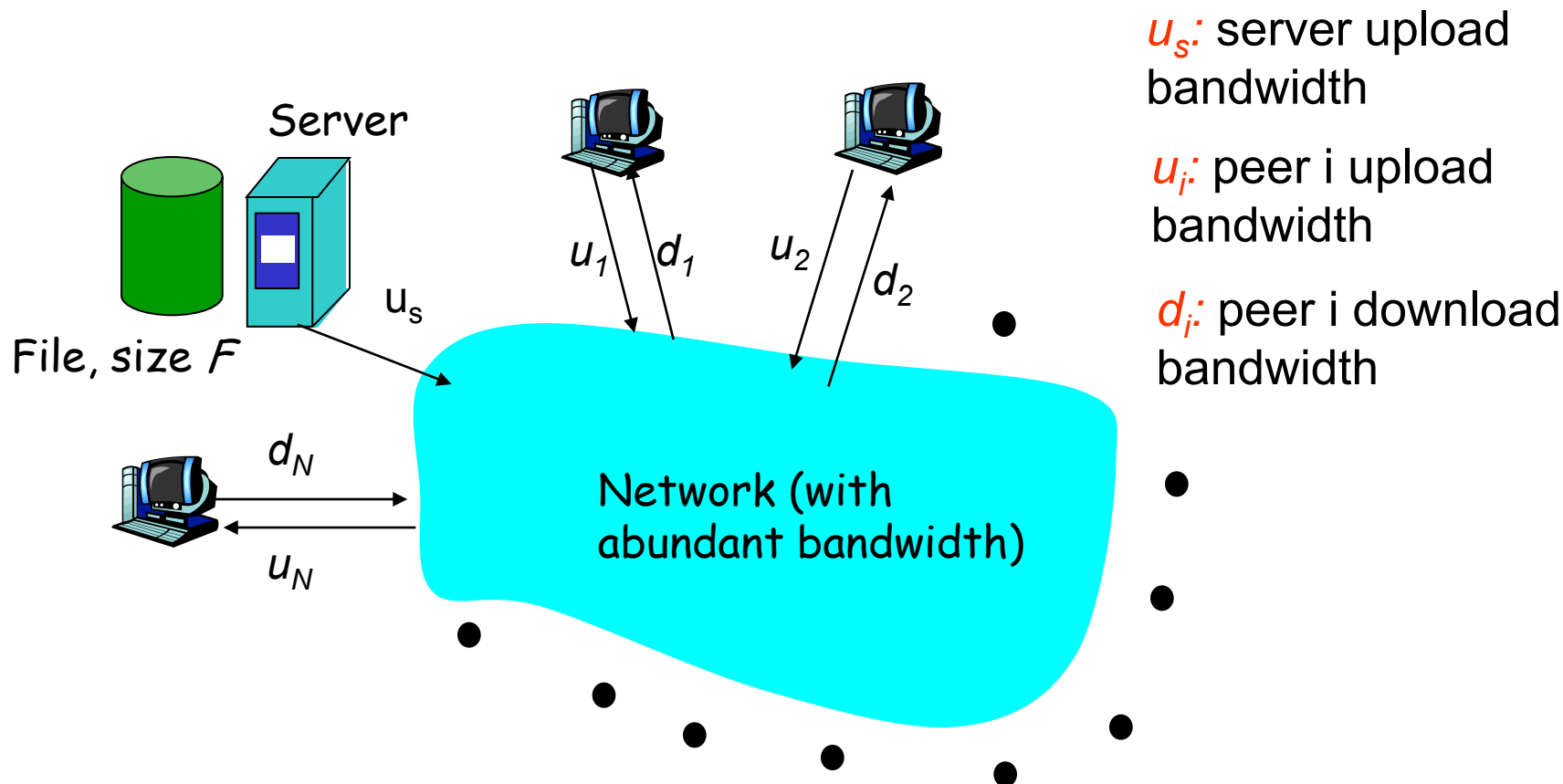
- ❑ Entire drives: C:\, D:\ (bad security risk)
- ❑ Just a folder or directory
- ❑ Printers
- ❑ USB flash drive
- ❑ Software and applications

Share Level Security

- ❑ Add the same accounts and passwords to all the computers on your network
- ❑ Each member can access the shared folder of the other computer, regardless of which computer they are using

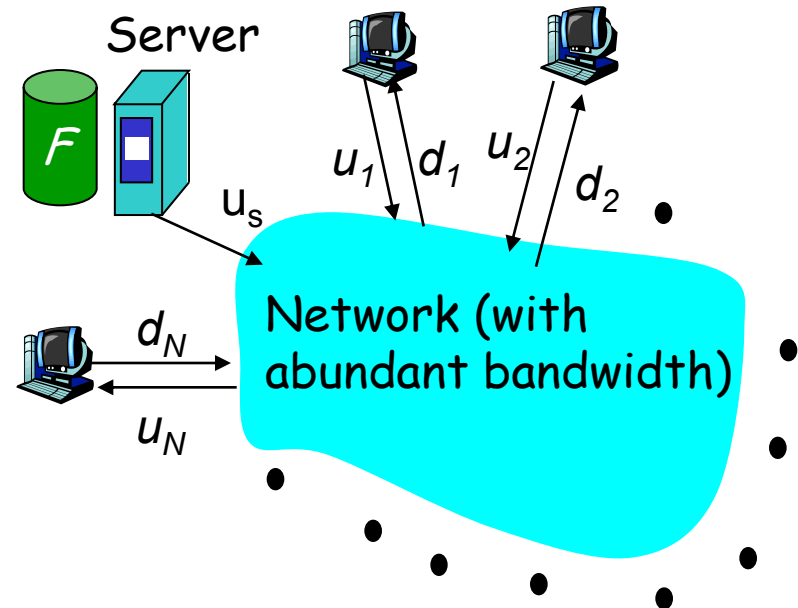
File Distribution: Server-Client vs P2P

Question: How much time to distribute file from one server to N peers?



File distribution time: server-client

- server sequentially sends N copies:
 - ❖ NF/u_s time
- client i takes F/d_i time to download

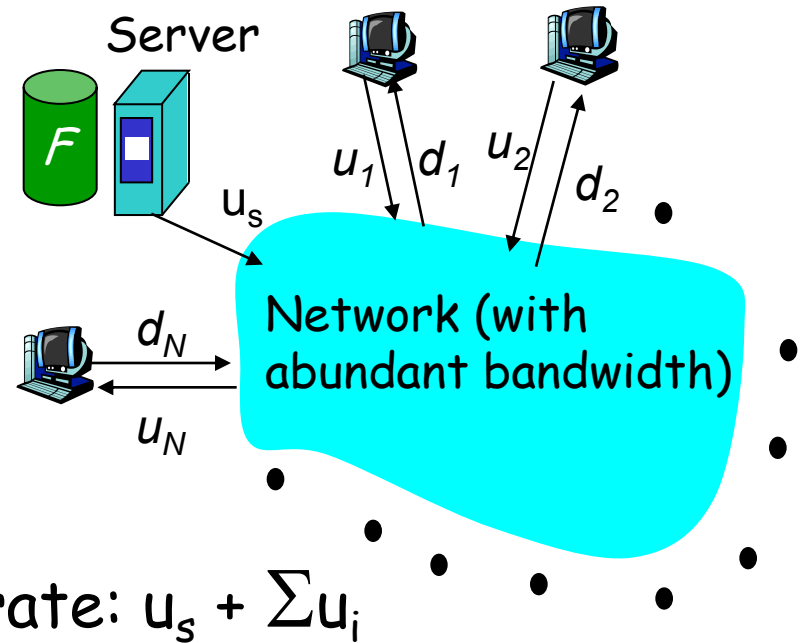


Time to distribute F
to N clients using client/server approach
 $= d_{cs} = \max \{ NF/u_s, F/\min_i(d_i) \}$

increases linearly in N
(for large N)

File distribution time: P2P

- ❑ server must send one copy: F/u_s time
- ❑ client i takes F/d_i time to download
- ❑ NF bits must be downloaded (aggregate)
 - ❑ fastest possible upload rate: $u_s + \sum u_i$



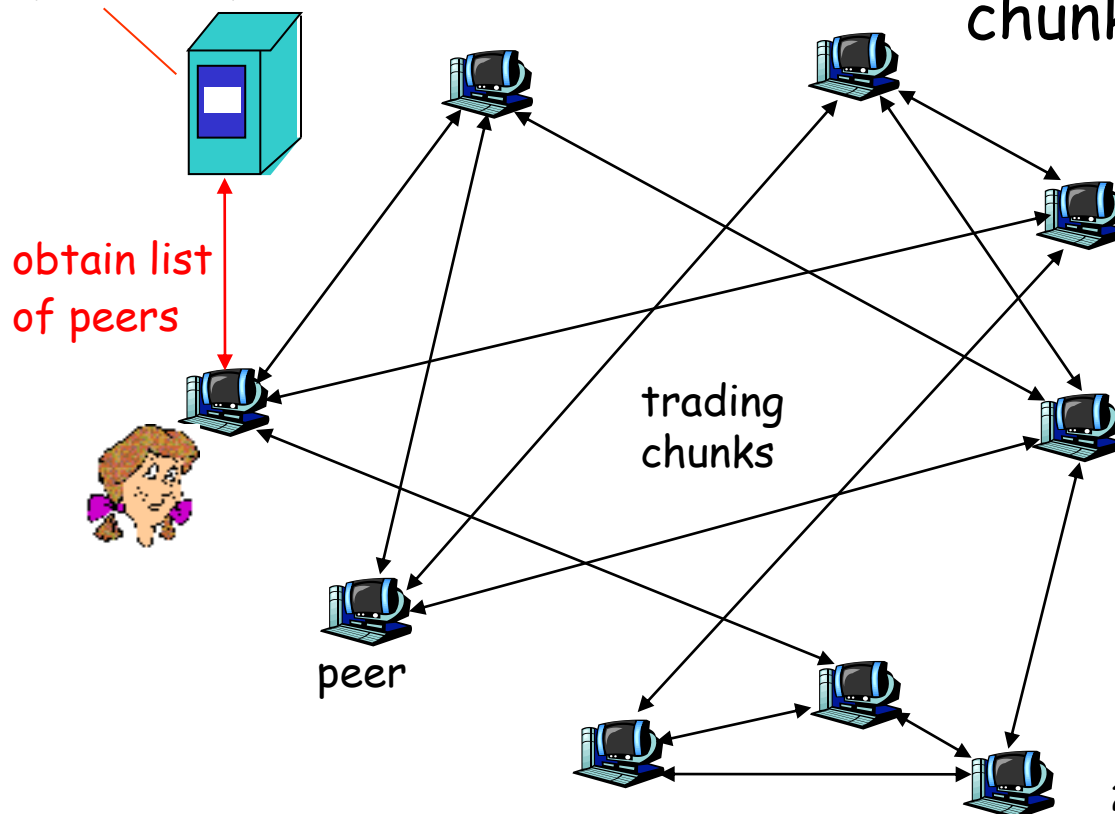
$$d_{p2p} = \max \left\{ F/u_s, F/\min(d_i)_i, NF/(u_s + \sum u_i) \right\}$$

File distribution: BitTorrent

□ P2P file distribution

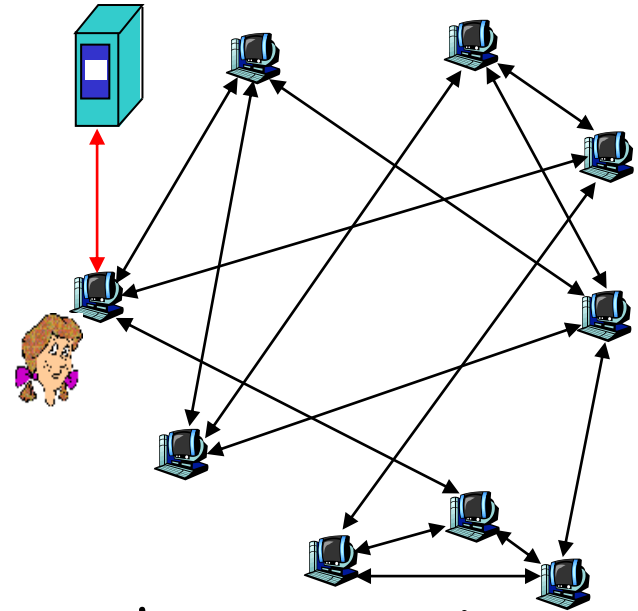
tracker: tracks peers participating in torrent

torrent: group of peers exchanging chunks of a file



BitTorrent (1)

- ❑ file divided into 256KB *chunks*.
- ❑ peer joining torrent:
 - ❖ has no chunks, but will accumulate them over time
 - ❖ registers with tracker to get list of peers, connects to subset of peers ("neighbors")
- ❑ while downloading, peer uploads chunks to other peers.
- ❑ peers may come and go
- ❑ once peer has entire file, it may (selfishly) leave or (altruistically) remain



BitTorrent (2)

Pulling Chunks

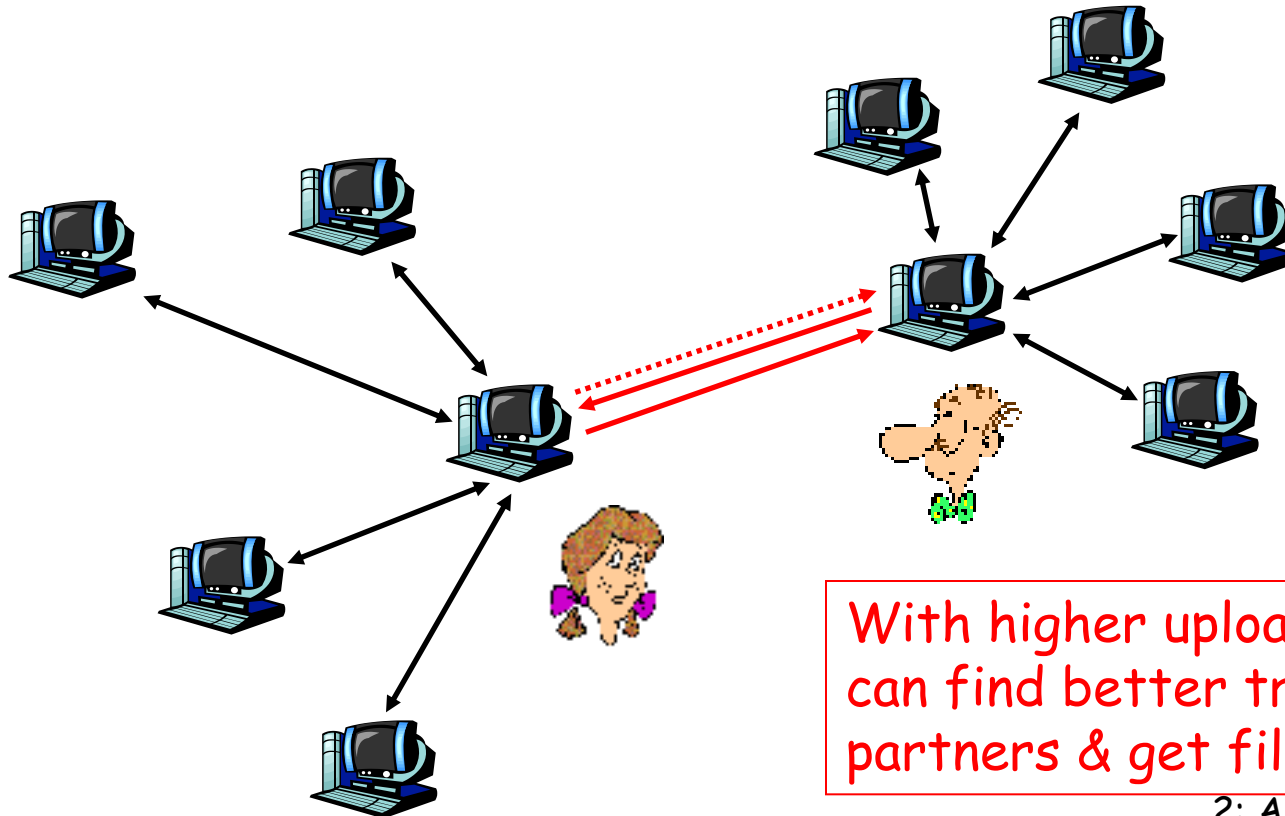
- at any given time, different peers have different subsets of file chunks
- periodically, a peer (Alice) asks each neighbor for list of chunks that they have.
- Alice sends requests for her missing chunks
 - ❖ rarest first

Sending Chunks: tit-for-tat

- Alice sends chunks to four neighbors currently sending her chunks *at the highest rate*
 - ❖ re-evaluate top 4 every 10 secs
- every 30 secs: randomly select another peer, starts sending chunks
 - ❖ newly chosen peer may join top 4
 - ❖ "optimistically unchoke"

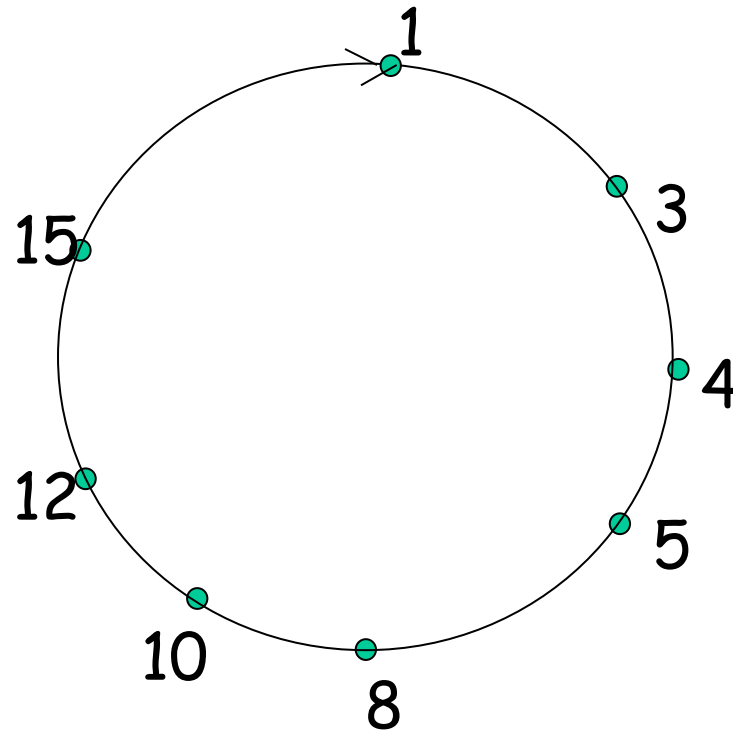
BitTorrent: Tit-for-tat

- (1) Alice "optimistically unchokes" Bob
- (2) Alice becomes one of Bob's top-four providers; Bob reciprocates
- (3) Bob becomes one of Alice's top-four providers



With higher upload rate,
can find better trading
partners & get file faster!

Circular DHT (1)



- ❑ Each peer *only* aware of immediate successor and predecessor.
- ❑ "Overlay network"

Distributed Hash Table (DHT)

- ❑ DHT = distributed P2P database
- ❑ Database has (key, value) pairs;
 - ❖ key: ss number; value: human name
 - ❖ key: content type; value: IP address
- ❑ Peers query DB with key
 - ❖ DB returns values that match the key
- ❑ Peers can also insert (key, value) peers

DHT Identifiers

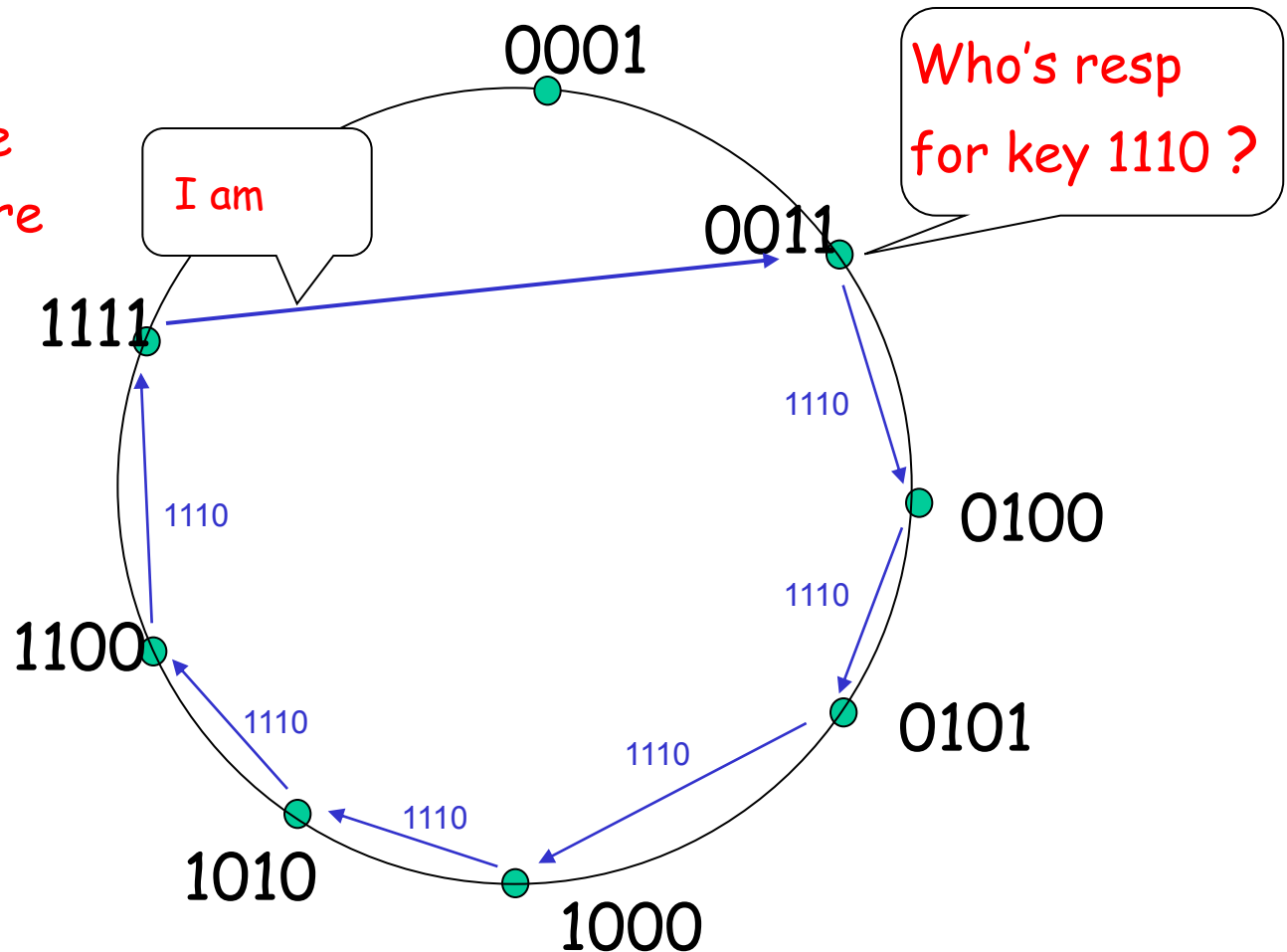
- ❑ Assign integer identifier to each peer in range $[0, 2^n - 1]$.
 - ❖ Each identifier can be represented by n bits.
- ❑ Require each key to be an integer in **same range**.
- ❑ To get integer keys, hash original key.
 - ❖ eg, $\text{key} = h(\text{"Led Zeppelin IV"})$
 - ❖ This is why they call it a distributed "hash" table

How to assign keys to peers?

- Central issue:
 - ❖ Assigning (key, value) pairs to peers.
- Rule: assign key to the peer that has the **closest** ID.
- Convention in lecture: closest is the **immediate successor** of the key.
- Ex: $n=4$; peers: 1,3,4,5,8,10,12,14;
 - ❖ key = 13, then successor peer = 14
 - ❖ key = 15, then successor peer = 1

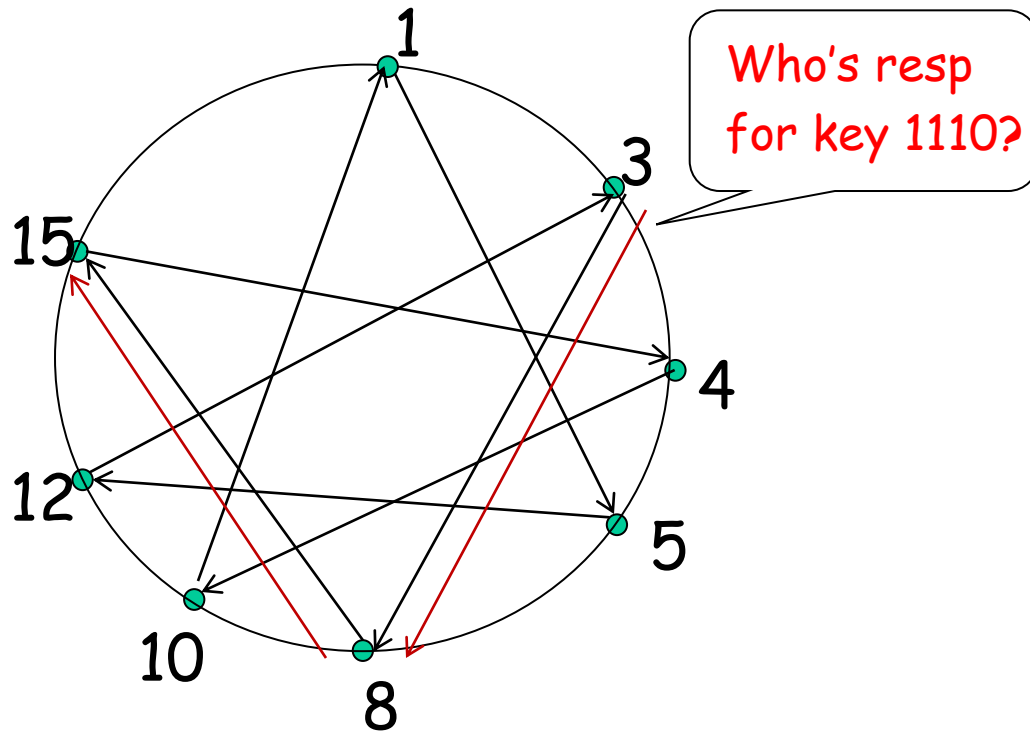
Circle DHT (2)

$O(N)$ messages
on avg to resolve
query, when there
are N peers



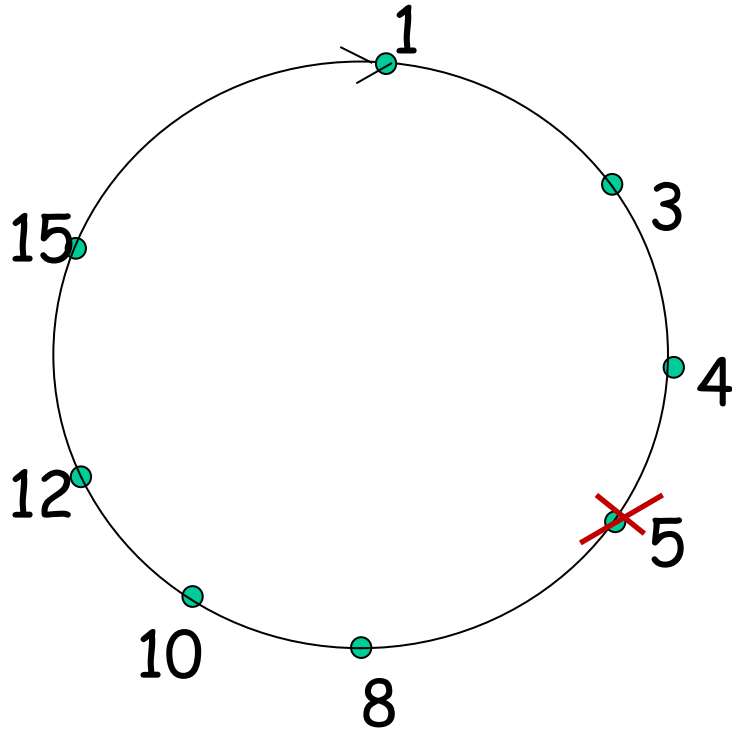
Define closest
as closest
successor

Circular DHT with Shortcuts



- ❑ Each peer keeps track of IP addresses of predecessor, successor, short cuts.
- ❑ Reduced from 6 to 2 messages.
- ❑ Possible to design shortcuts so $O(\log N)$ neighbors, $O(\log N)$ messages in query

Peer Churn

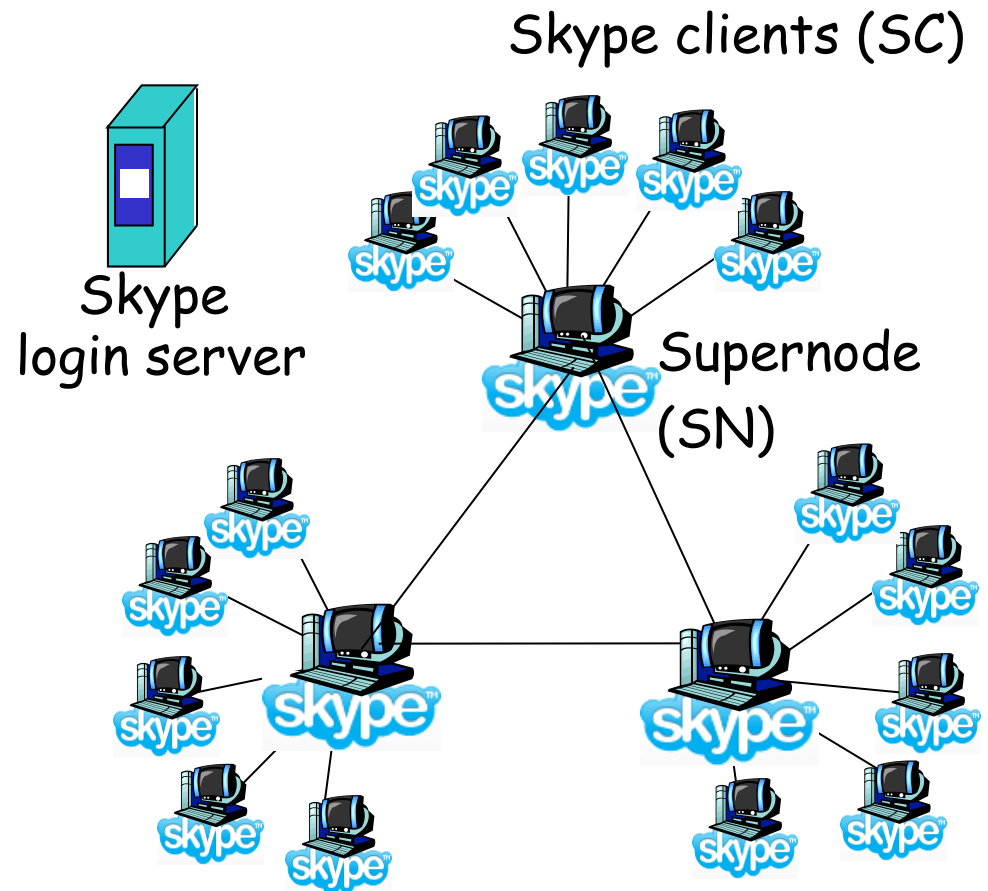


- To handle peer churn, require each peer to know the IP address of its two successors.
- Each peer periodically pings its two successors to see if they are still alive.

- ❑ Peer 5 abruptly leaves
- ❑ Peer 4 detects; makes 8 its immediate successor; asks 8 who its immediate successor is; makes 8's immediate successor its second successor.
- ❑ What if peer 13 wants to join?

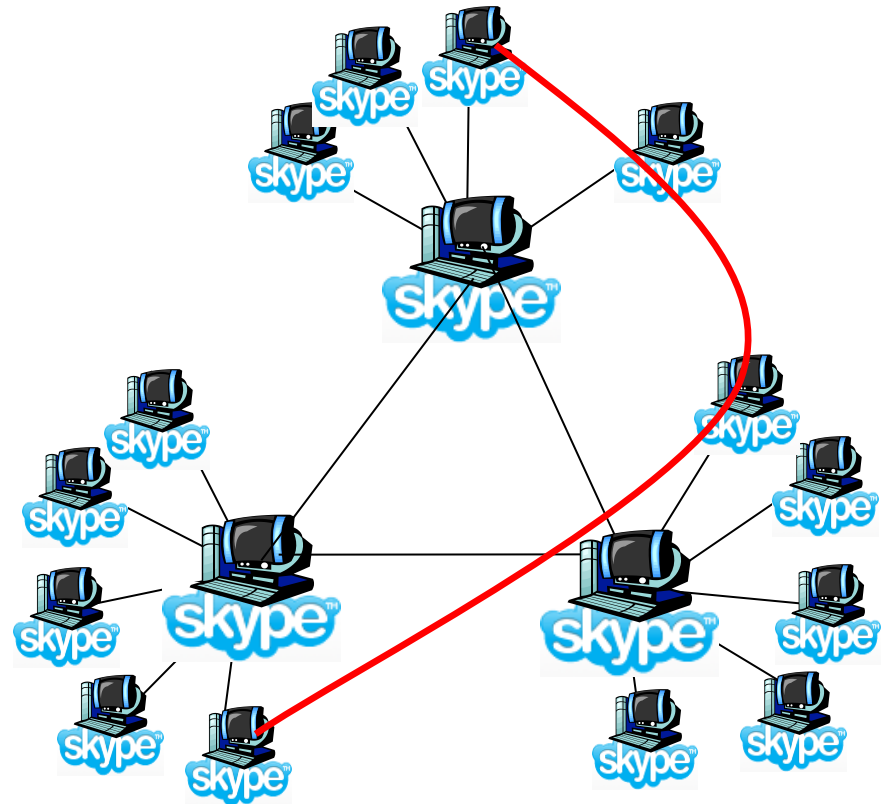
P2P Case study: Skype

- ❑ inherently P2P: pairs of users communicate.
- ❑ proprietary application-layer protocol (inferred via reverse engineering)
- ❑ hierarchical overlay with SNs
- ❑ Index maps usernames to IP addresses; distributed over SNs



Peers as relays

- ❑ Problem when both Alice and Bob are behind "NATs".
 - ❖ NAT prevents an outside peer from initiating a call to insider peer
- ❑ Solution:
 - ❖ Using Alice's and Bob's SNs, Relay is chosen
 - ❖ Each peer initiates session with relay.
 - ❖ Peers can now communicate through NATs via relay



Chapter 2: Summary

our study of network apps now complete!

- application architectures
 - ❖ client-server
 - ❖ P2P
 - ❖ hybrid
- application service requirements:
 - ❖ reliability, bandwidth, delay
- Internet transport service model
 - ❖ connection-oriented, reliable: TCP
 - ❖ unreliable, datagrams: UDP
- specific protocols:
 - ❖ HTTP
 - ❖ FTP
 - ❖ SMTP, POP, IMAP
 - ❖ DNS
 - ❖ P2P: BitTorrent, Skype
- socket programming

Chapter 2: Summary

Most importantly: learned about *protocols*

□ typical request/reply message exchange:

- ❖ client requests info or service
- ❖ server responds with data, status code

□ message formats:

- ❖ headers: fields giving info about data
- ❖ data: info being communicated

Important themes:

- control vs. data msgs
 - ❖ in-band, out-of-band
- centralized vs. decentralized
- stateless vs. stateful
- reliable vs. unreliable msg transfer
- "complexity at network edge"